

# Logic

*Logic* is a discipline that studies the principles and methods used in correct reasoning

It includes:

- A formal language for expressing statements.
- An inference mechanism (a collection of rules) to reason about valid arguments.

# Logic

*Logic* is a discipline that studies the principles and methods used to construct valid arguments.

- ◆ An *argument* is a related sequence of statements to demonstrate the truth of an assertion
  - *premises* are assumed to be true
  - *conclusion*, the last statement of the sequence, is taken to be true based on the truth of the other statements.
  
- ◆ An *argument* is valid if the conclusion follows logically from the truth of the premises.
  
- ◆ *Logic* is the foundation for expressing formal proofs.

# Propositional Logic

*Propositional Logic* is the logic of compound statements built from simpler statements using *Logical Boolean connectives*.

## Some applications

- Design of digital electronic circuits.
- Expressing conditions in programs.
- Queries to databases & search engines.

# Propositional Logic

A *proposition* is a declarative sentence that is either TRUE or FALSE (not both).

Examples:

- The Earth is flat
- $3 + 2 = 5$
- I am older than my mother
- Tallahassee is the capital of Florida
- $5 + 3 = 9$
- Athens is the capital of Georgia

# Propositional Logic

- ❖ Letters are used to denote propositions.
  - ❖ The most frequently used letters are  $p$ ,  $q$ ,  $r$ ,  $s$  and  $t$ .
- ❖ Example:

$p$ : Grass is green.

- ❖ **The truth value** of a proposition is true, denoted by  $T$ , if it is true, and false, denoted by  $F$  if it is false

# Propositional Logic

- ◆ **Compound propositions** : built up from simpler propositions using logical operators
  - Frequently corresponds with compound English sentences.

## Example:

Given

p: Jack is older than Jill  
q: Jill is female

We can build up

r: Jack is older than Jill and Jill is female ( $p \wedge q$ )

s: Jack is older than Jill or Jill is female ( $p \vee q$ )

t: Jack is older than Jill and it is not the case that Jill is female  
( $p \wedge \neg q$ )

# Propositional Logic - negation

Let  $p$  be a proposition.

The *negation* of  $p$  is written  $\neg p$  and has meaning:

"It is not the case that  $p$ ."

Truth table for negation:

$p$	$\neg p$
T	F
F	T

# Propositional Logic - conjunction

- ◆ *Conjunction operator* " $\wedge$ " (*AND*):
  - corresponds to English "and."
  - is a *binary* operator in that it operates on two propositions when creating a compound proposition
- ◆ Def. Let  $p$  and  $q$  be two arbitrary propositions, the conjunction of  $p$  and  $q$ , denoted

$$p \wedge q,$$

*is true* if both  $p$  and  $q$  are *true* and *false* otherwise.

# Propositional Logic - conjunction

Conjunction operator

$p \wedge q$  is true when  $p$  and  $q$  are both true.

Truth table for conjunction:

$p$	$q$	$p \wedge q$
T	T	T
T	F	F
F	T	F
F	F	F

# Propositional Logic - disjunction

Disjunction operator  $\vee$  (or):

- **loosely** corresponds to English "or."
- binary operator

Def.: Let  $p$  and  $q$  be two arbitrary propositions, the **disjunction** of  $p$  and  $q$ , denoted

$$p \vee q$$

is false when both  $p$  and  $q$  are false and true otherwise.



$\vee$  is also called **inclusive or**

- Observe that  $p \vee q$  is true when  $p$  is true, or  $q$  is true, or both  $p$  and  $q$  are true.

# Propositional Logic - disjunction

Disjunction operator

$p \vee q$  is true when  $p$  or  $q$  (or both) is true.

Truth table for conjunction:

$p$	$q$	$p \vee q$
T	T	T
T	F	T
F	T	T
F	F	F

# Propositional Logic - XOR

Exclusive Or operator ( $\oplus$ ):

- ❖ corresponds to English "either...or..." (exclusive form of or)
- ❖ binary operator

Def.: Let  $p$  and  $q$  be two arbitrary propositions, the exclusive or of  $p$  and  $q$ , denoted

$$p \oplus q$$

is true when either  $p$  or  $q$  (but not both) is true.

# Propositional Logic - XOR

Exclusive Or:

$p \oplus q$  is true when  $p$  or  $q$  (not both) is true.

Truth table for exclusive or:

$p$	$q$	$p \oplus q$
T	T	F
T	F	T
F	T	T
F	F	F

# Propositional Logic- Implication

## Implication operator ( $\rightarrow$ ):

- ❖ binary operator
- ❖ similar to the English usage of "if...then...", "implies", and many other English phrases

**Def.:** Let  $p$  and  $q$  be two arbitrary propositions, the implication  $p \rightarrow q$  is false when  $p$  is false and  $q$  is true, and true otherwise.

$p \rightarrow q$  is true when  $p$  is true and  $q$  is true,  $q$  is true, or  $p$  is false.

$p \rightarrow q$  is false when  $p$  is true and  $q$  is false.

Example:

$r$ : "The dog is barking."

$s$ : "The dog is awake."

$r \rightarrow s$ : "If the dog is barking then the dog is awake."

# Propositional Logic- Implication

Truth table for implication:

$p$	$q$	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

- ❖ If the temperature is below  $10^{\circ}$  F, then water freezes.

# Propositional Logic- Biconditional

## Biconditional operator ( $\leftrightarrow$ ):

- ❖ Binary operator
- ❖ Partly similar to the English usage of "If and only if"

Def.: Let  $p$  and  $q$  be two arbitrary propositions.

The *biconditional*  $p \leftrightarrow q$  is true when  $q$  and  $p$  have the same truth values and false otherwise.

Example:

$p$ : "The dog plays fetch."

$q$ : "The dog is outside."

$p \leftrightarrow q$ : "The plays fetch if and only if it is outside."

# Propositional Logic- Biconditional

Truth table for biconditional:

$p$	$q$	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

# Propositional Equivalences

◆ A *tautology* is a proposition that is always true.

■ Ex.:  $p \vee \neg p$

$p$	$\neg p$	$p \vee \neg p$
T	F	T
F	T	T

◆ A *contradiction* is a proposition that is always false.

■ Ex.:  $p \wedge \neg p$

$p$	$\neg p$	$p \wedge \neg p$
T	F	F
F	T	F

◆ A *contingency* is a proposition that is neither a tautology nor a contradiction.

■ Ex.:  $p \rightarrow \neg p$

$p$	$\neg p$	$p \rightarrow \neg p$
T	F	F
F	T	T

# Propositional Logic: Logical Equivalence

- ◆ If  $p$  and  $q$  are propositions, then  $p$  is **logically equivalent to  $q$**  if their truth tables are the same.
  - " $p$  is equivalent to  $q$ ." is denoted by  $p \equiv q$
- ◆  $p, q$  are *logically equivalent* if their biconditional  $p \leftrightarrow q$  is a tautology.

# Propositional Logic: Logical Equivalence

$$\diamond p \equiv \neg \neg p$$

$p$	$\neg p$	$\neg \neg p$
T	F	T
F	T	F

The equivalence holds since these two columns are the same.

# Propositional Logic: Logical Equivalence

◆  $p \rightarrow q \equiv \neg p \vee q$ ?

◆ Does  $(p \rightarrow q) \leftrightarrow (\neg p \vee q)$ ?

◆ Does  $(p \rightarrow q) \rightarrow (\neg p \vee q)$  and  
 $(\neg p \vee q) \rightarrow (p \rightarrow q)$ ?

# Propositional Logic: Logical Equivalences

- *Identity*

$$p \wedge \mathbf{T} \equiv p$$

$$p \vee \mathbf{F} \equiv p$$

- *Domination*

$$p \vee \mathbf{T} \equiv \mathbf{T}$$

$$p \wedge \mathbf{F} \equiv \mathbf{F}$$

- *Idempotent*

$$p \vee p \equiv p$$

$$p \wedge p \equiv p$$

- *Double negation*

$$\neg(\neg p) \equiv p$$

# Propositional Logic: Logical Equivalences

- *Commutative:*

$$p \vee q \equiv q \vee p$$

$$p \wedge q \equiv q \wedge p$$

- *Associative:*

$$(p \vee q) \vee r \equiv p \vee (q \vee r)$$

$$(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$$

# Propositional Logic: Logical Equivalences

- *Distributive:*

$$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$$

$$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$$

- *De Morgan's:*

$$\neg(p \wedge q) \equiv \neg p \vee \neg q \quad (\text{De Morgan's I})$$

$$\neg(p \vee q) \equiv \neg p \wedge \neg q \quad (\text{De Morgan's II})$$

# Propositional Logic: Logical Equivalences

- *Absorption:*

$$p \vee (p \wedge q) \equiv p$$

$$p \wedge (p \vee q) \equiv p$$

- *Negation:*

$$p \wedge \neg p \equiv \mathbf{F}$$

$$p \vee \neg p \equiv \mathbf{T}$$

- A useful LE involving  $\rightarrow$ :

$$p \rightarrow q \equiv \neg p \vee q$$

# Predicate Logic

Define:

**UGA**( $x$ ) = "x is a UGA student."

**Universe of Discourse** – *all people*

$x$  is a variable that represents an arbitrary individual in the Universe of Discourse

A **predicate**  $P$ , or propositional function, is a function that maps objects in the universe of discourse to propositions

- **UGA**(Paris Hilton) is a **proposition**.
- **UGA**( $x$ ) is **not a proposition**.

UGA( $x$ ) is like an English predicate template

- \_\_\_\_\_ is a UGA student

# Predicate Logic

◆  $\text{Pos}(x): x > 0$

**Universe of Discourse (UoD): Integers**

$\text{Pos}(1): 1 > 0 = \text{T}$

$\text{Pos}(50): 50 > 0 = \text{T}$

$\text{Pos}(-10): -10 > 0 = \text{F}$

◆  $\text{Female}(x): x \text{ is female}$

UoD: all people

$\text{Female}(\text{Maria}) = \text{T}$

$\text{Female}(\text{Edward}) = \text{F}$

# Predicate Logic

- ◆ A predicate that states a property about one object is called a monadic predicate.
  - UGA-Student(x)
  - Parent(x)
  - Female(x)

# Predicate Logic

- ◆ A predicate of the form  $P(x_1, \dots, x_n)$ ,  $n > 1$  that states the relationships among the objects  $x_1, \dots, x_n$  is called polyadic.

Also, an n-place predicate or n-ary predicate (a predicate with arity n).

- ◆ It takes  $n > 1$  arguments
- ◆  $UoD(x_1, \dots, x_n) = UoD(x_1) \times UoD(x_2) \times \dots \times UoD(x_n)$

- ◆  $L(x, y)$ : x loves y

$UoD(x) = UoD(y) = \text{all people}$

$L(\text{Adam}, \text{Mary}) = T$

$L(\text{Mike}, \text{Jill}) = F$

- ◆  $Q(x, y, z)$ :  $x + y = z$

$UoD(x) = UoD(y) = UoD(z)$ : positive integers

$Q(1, 1, 2) = T$

$Q(1, 2, 5) = F$

# Predicate Logic: Universal Quantifier

Suppose that  $P(x)$  is a predicate on some universe of discourse.

The universal quantification of  $P(x)$  ( $\forall x P(x)$ ) is the **proposition**:

" $P(x)$  is true for all  $x$  in the universe of discourse."

$\forall x P(x)$  reads "for all  $x$ ,  $P(x)$  is True"

- ◆  $\forall x P(x)$  is TRUE means  $P(x)$  is true for all  $x$  in  $UoD(x)$ .
- ◆  $\forall x P(x)$  is FALSE means there is an  $x$  in  $UoD(x)$  for which  $P(x)$  is false.

# Predicate Logic: Universal Quantifier

◆  $ID(x)$  means  $x$  has a student ID.

$UoD(x)$ : UGA students

$\forall x ID(x)$  means

every UGA student has a student ID

◆  $IsFemale(x)$  means  $x$  is a Female

$UoD(x)$ : UGA students

$\forall x IsFemale(x)$  means

Every UGA student is a female!

Jim is a UGA student who is a male

Jim is called a **counterexample** for  $\forall x IsFemale(x)$

# Predicate Logic: Universal Quantifier

In the special case that the universe of discourse  $U$ , is finite, ( $U = \{a_1, a_2, a_3, \dots, a_n\}$ )

$$\forall x P(x)$$

corresponds to the proposition:

$$P(a_1) \wedge P(a_2) \wedge \dots \wedge P(a_n)$$

We can write a program to loop through the elements in the universe and check each for truthfulness.

If all are true, then the proposition is true.

Otherwise it is false!

# Predicate Logic: Existential Quantifier

Suppose  $P(x)$  is a predicate on some universe of discourse.

The existential quantification of  $P(x)$  is the proposition:

“There exists at least one  $x$  in the universe of discourse such that  $P(x)$  is true.”

$\exists x P(x)$  reads “for some  $x$ ,  $P(x)$ ” or “There exists  $x$ ,  $P(x)$  is True”

$\exists x P(x)$  is **TRUE** means  
there is an  $x$  in  $UoD(x)$  for which  $P(x)$  is true.

$\exists x P(x)$  is **FALSE** means :  
for all  $x$  in  $UoD(x)$ ,  $P(x)$  is false

# Predicate Logic: Existential Quantifier

Examples:

◆  $F(x)$  means  $x$  is female.

UoD( $x$ ): UGA students

$\exists x F(x)$  means

- For some UGA student  $x$ ,  $x$  is a female.
- There exists a UGA student  $x$  who is a female.

◆  $Y(x)$  means  $x$  is less than 13 years old

UoD( $x$ ): UGA students

$\exists x Y(x)$  means

- ◆ for some UGA student  $x$ ,  $x$  is less than 13 years old
- ◆ there exists a UGA student  $x$  such that  $x$  is less than 13 years old

What will be the truth-value of  $\exists x Y(x)$  if the UoD is all people?

# Predicate Logic: Existential Quantifier

In the special case that the universe of discourse,  $U$ , is finite, ( $U = \{x_1, x_2, x_3, \dots, x_n\}$ )

$$\exists x P(x)$$

corresponds to the proposition:

$$P(x_1) \vee P(x_2) \vee \dots \vee P(x_n)$$

We can write a program to loop through the elements in the universe and check each for truthfulness. If all are false, then the proposition is false. Otherwise, it is true!

# Predicates - Quantifier negation

$\exists x P(x)$  means "P(x) is true for some x."

What about  $\neg \exists x P(x)$  ?

It is not the case that ["P(x) is true for some x."]

"P(x) is not true for all x."

$$\forall x \neg P(x)$$

Existential negation:

$$\neg \exists x P(x) \equiv \forall x \neg P(x).$$

# Predicates - Quantifier negation

$\forall x P(x)$  means "P(x) is true for every x."

What about  $\neg \forall x P(x)$  ?

It is not the case that ["P(x) is true for every x."]

"There exists an x for which P(x) is not true."

$$\exists x \neg P(x)$$

Universal negation:

$$\neg \forall x P(x) \equiv \exists x \neg P(x).$$

# Re-Cap

A **predicate**  $P$ , or propositional function, is a function that maps objects in the **universe of discourse** to propositions

- ◆ Predicates can be quantified using the universal quantifier (“for all”)  $\forall$  or the existential quantifier (“there exists”)  $\exists$
- ◆ Quantified predicates can be negated as follows
  - $\neg \forall x P(x) \equiv \exists x \neg P(x)$
  - $\neg \exists x P(x) \equiv \forall x \neg P(x)$
- ◆ Quantified variables are called “bound”
- ◆ Variables that are not quantified are called “free”

# Proofs

A *theorem* is a statement that can be proved to be true.

A *proof* is a sequence of statements that form an argument.

# Proofs: Inference Rules

◆ An *Inference Rule*:

$$\begin{array}{l} \textit{premise 1} \\ \textit{premise 2 ...} \\ \hline \therefore \textit{conclusion} \end{array}$$

" $\therefore$ " means "therefore"

# Proofs: Modus Ponens

I have a total score over 96.

If I have a total score over 96, then I get an A for the class.

$\therefore$  I get an A for this class

$$\begin{array}{c} p \\ p \rightarrow q \\ \hline \therefore q \end{array}$$

Tautology: $(p \wedge (p \rightarrow q)) \rightarrow q$
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# Proofs: Modus Tollens

If the power supply fails then the lights go out.

The lights are on.

$\therefore$  The power supply has not failed.

$$\begin{array}{c} \neg q \\ p \rightarrow q \\ \hline \therefore \neg p \end{array}$$

Tautology:

$$(\neg q \wedge (p \rightarrow q)) \rightarrow \neg p$$

# Proofs: Addition

I am a student.

$\therefore$  I am a student or I am a visitor.

$$\frac{p}{\therefore p \vee q}$$

Tautology: $p \rightarrow (p \vee q)$
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# Proofs: Simplification

I am a student and I am a soccer player.

$\therefore$  I am a student.

$$\frac{p \wedge q}{\therefore p}$$

Tautology:

$$(p \wedge q) \rightarrow p$$

# Proofs: Conjunction

I am a student.

I am a soccer player.

$\therefore$  I am a student and I am a soccer player.

$$\frac{p}{q}$$

---

$$\therefore p \wedge q$$

Tautology:

$$((p) \wedge (q)) \rightarrow p \wedge q$$

# Proofs: Disjunctive Syllogism

I am a student or I am a soccer player.

I am a not soccer player.

$\therefore$  I am a student.

$$\begin{array}{c} p \vee q \\ \neg q \\ \hline \therefore p \end{array}$$

Tautology:

$$((p \vee q) \wedge \neg q) \rightarrow p$$

# Proofs: Hypothetical Syllogism

If I get a total score over 96, I will get an A in the course.

If I get an A in the course, I will have a 4.0 average.

∴ If I get a total score over 96 then I will have a 4.0 average.

$$\begin{array}{l} p \rightarrow q \\ q \rightarrow r \\ \hline \therefore p \rightarrow r \end{array}$$

Tautology:

$$((p \rightarrow q) \wedge (q \rightarrow r)) \rightarrow (p \rightarrow r)$$

# Proofs: Resolution

I am taking CS4540 or I am taking CS4560.

I am not taking CS4540 or I am taking CS7300.

$\therefore$  I am taking CS4560 or I am taking CS7300.

$$\frac{p \vee q \quad \neg p \vee r}{\therefore q \vee r}$$

Tautology:

$$((p \vee q) \wedge (\neg p \vee r)) \rightarrow (q \vee r)$$

# Fallacy of Affirming the Conclusion

If you have the flu then you'll have a sore throat.

You have a sore throat.

∴ You must have the flu.

$$\frac{q \quad p \rightarrow q}{\therefore p}$$

Fallacy:

$$(q \wedge (p \rightarrow q)) \rightarrow p$$

# Fallacy of Denying the Hypothesis

If you have the flu then you'll have a sore throat.

You do not have the flu.

∴ You do not have a sore throat.

$$\frac{\begin{array}{l} \neg p \\ p \rightarrow q \end{array}}{\therefore \neg q}$$

Fallacy:

$$(\neg p \wedge (p \rightarrow q)) \rightarrow \neg q$$

# Inference Rules for Quantified Statements

$$\frac{\forall x P(x)}{\therefore P(c)}$$

## Universal Instantiation

(for an arbitrary object  $c$  from UoD)

$$\frac{P(c)}{\therefore \forall x P(x)}$$

## Universal Generalization

(for any arbitrary element  $c$  from UoD)

$$\frac{\exists x P(x)}{\therefore P(c)}$$

## Existential Instantiation

(for some specific object  $c$  from UoD)

$$\frac{P(c)}{\therefore \exists x P(x)}$$

## Existential Generalization

(for some object  $c$  from UoD)